



The isotopy problem of braids

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- The braid isotopy problem is a problem of **medium** difficulty, with **many** (really) different solutions illustrating various approaches to Artin's braid groups.
- Here: a survey of **some** solutions:
 - ▶ one **algebraic** solution: the greedy normal form
 - ▶ two **topological** solutions: Dynnikov's coordinates, Bressaud's relaxation method [and two more: the alternating normal form (yesterday), handle reduction (ILDT)]

Plan:

- 1. The braid isotopy problem
- 2. Greedy normal form and the Garside structure
- 3. Dynnikov's coordinates
- 4. Bressaud's relaxation algorithm

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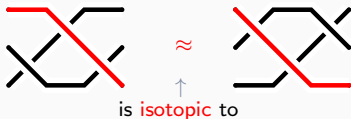
- A 3-strand **braid diagram**:



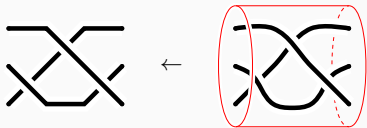
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- **Isotopy** Problem:

Given two n -strand braid diagrams, can one **deform** them to one another?

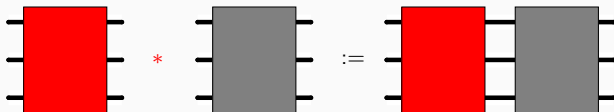


- More formally: view braid diagrams as projections of 3D-diagrams in $D^2 \times (0, 1)$,

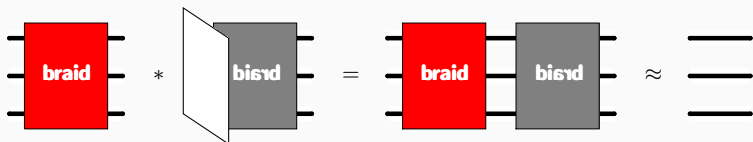


and consider **ambient isotopy** leaving the end-disks fixed.

- **Concatenation** of braid diagrams:



- ▶ Associative;
- ▶ Compatible with isotopy, hence induces a well-defined product on classes;
- ▶ Admits the unbraided diagram $[\emptyset]$ as a neutral element;
- ▶ Every diagram has an inverse, its **mirror-image**:



- For every $n \geq 1$: the **group** B_n of n -strand **braids**.

↑
isotopy class of braid diagrams

- The group structure of B_n makes the Braid Isotopy Problem easier:
 - Reduces to the Braid **Triviality** Problem: $D' \approx D \Leftrightarrow D^{-1} * D' \approx [\emptyset]$.
 - Enables one to use algebraic tools, provided one has a **presentation** of B_n .
- Artin** generators: Every n -strand braid diagram is a (**finite**) concatenation of elementary diagrams with **one** crossing, hence of the form

$$\sigma_i : \begin{array}{c} \text{---} n \\ \vdots \\ \text{---} \\ \text{---} \\ \text{---} \\ \vdots \\ \text{---} \\ \vdots \\ \text{---} \\ \vdots \\ \text{---} 1 \end{array} \begin{array}{c} i+1 \\ \diagdown \\ \diagup \\ i \end{array} \quad \text{or} \quad \sigma_i^{-1} : \begin{array}{c} \text{---} n \\ \vdots \\ \text{---} \\ \text{---} \\ \text{---} \\ \vdots \\ \text{---} \\ \vdots \\ \text{---} \\ \vdots \\ \text{---} 1 \end{array} \begin{array}{c} i+1 \\ \diagup \\ \diagdown \\ i \end{array} \quad \text{with } 1 \leq i < n.$$

- Theorem** (Artin, 1926): The group B_n admits the presentation

$$\langle \sigma_1, \dots, \sigma_{n-1} \mid \begin{array}{l} \sigma_i \sigma_j = \sigma_j \sigma_i \quad \text{for } |i-j| \geq 2 \\ \sigma_i \sigma_j \sigma_i = \sigma_j \sigma_i \sigma_j \quad \text{for } |i-j| = 1 \end{array} \rangle.$$

► Proof: Isotopy of piecewise linear diagrams is generated by Δ -moves. □

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- Braid Isotopy reduced to the **Word Problem** for B_n with respect to $\{\sigma_1, \dots, \sigma_{n-1}\}$:
given a **braided word** w , decide whether w represents 1 in B_n .

↑
a word in the letters $\sigma_1^{\pm 1}, \dots, \sigma_{n-1}^{\pm 1}$.

- (Novikov, 1952) There exists a finitely presented group with an unsolvable Word Problem.
- Here: (Garside) Use the **monoid**.

• Theorem (Garside, 1969): Let B_n^+ be the monoid with presentation

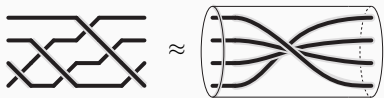
$$\left\langle \sigma_1, \dots, \sigma_{n-1} \mid \begin{array}{ll} \sigma_i \sigma_j = \sigma_j \sigma_i & \text{for } |i-j| \geq 2 \\ \sigma_i \sigma_j \sigma_i = \sigma_j \sigma_i \sigma_j & \text{for } |i-j| = 1 \end{array} \right\rangle^+.$$

Then B_n^+ embeds in B_n and B_n is a **group of fractions** for B_n^+ .

↑
every element of B_n can be written $\beta^{-1}\gamma$ with $\beta, \gamma \in B_n^+$

► Proof: Show that B_n^+ is cancellative and admits common multiples. □

- An **effective** way of reducing from B_n to B_n^+ :
- Lemma (Garside): *Inductively define Δ_n by $\Delta_1 = 1$, $\Delta_n = \Delta_{n-1} \cdot \sigma_{n-1} \cdots \sigma_2 \sigma_1$.*



Then, for every (signed) n -strand braid word w , one can find $p \geq 0$
and a positive n -strand braid word w' and satisfying $\Delta_n^p w \equiv w'$.

- Then: $w \equiv \epsilon \Leftrightarrow w' \equiv \Delta_n^p \Leftrightarrow w' \equiv^+ \Delta_n^p$
 ↑ ↑ ↑
 the empty word equivalence equivalence
 generated by braid relations generated by braid relations alone
 and $\sigma_i \sigma_i^{-1} = \sigma_i^{-1} \sigma_i = 1$

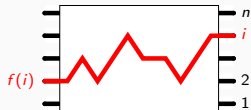
- Now: \equiv^+ is decidable, as it preserves **word-length**.
- Hence: A (theoretical) solution to the Braid Isotopy Problem: starting from w ,
 - ▶ 1. find p and w' positive satisfying $\Delta_n^p w \equiv w'$;
 - ▶ 2. test $w' \equiv^+ \Delta_n^p$ by systematically enumerating the \equiv^+ -class of w' .

- To improve the previous solution and make it tractable:
define (efficiently computable) **normal forms** on B_n^+ .

- Every n -strand braid gives a permutation of $\{1, \dots, n\}$:
follow the positions of the strands:

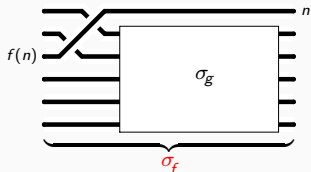
- ▶ short exact sequence

$$1 \longrightarrow PB_n \longrightarrow B_n \longrightarrow \mathfrak{S}_n \longrightarrow 1.$$



- Inductively define a (set-theoretic) **section**
for the projection of B_n onto \mathfrak{S}_n :
for $f = (n, f(n)) \circ g$ with $g \in \mathfrak{S}_{n-1}$,
put $\sigma_f := \sigma_{f(n)} \cdots \sigma_{n-1} \sigma_g$

- ▶ a family of $n!$ **permutation braids** in B_n^+ .



- Lemma: Permutations braids are the (left- and right-) **divisors** of Δ_n in B_n^+ .

β left-divides γ if $\exists \gamma' (\beta\gamma' = \gamma)$.

- Theorem (Garside 1969): *With respect to (left- and right-) divisibility, B_n^+ is a **lattice**.*

↑
least common multiples and greatest common divisors exist

- Corollary: *For every positive n -strand braid β , there exists a **unique maximal** permutation braid left-dividing β .*

↑
namely: the left-gcd of β and Δ_n

- ▶ A distinguished decomposition:

$$\beta = \sigma_{f_1} \cdot \beta' = \sigma_{f_1} \cdot \sigma_{f_2} \cdot \beta'' = \dots = \sigma_{f_1} \cdot \sigma_{f_2} \cdot \dots \cdot \sigma_{f_r}.$$

“a positive braid is a sequence of permutations”

- Fact: σ_f is a maximal left-divisor of $\sigma_f \cdot \sigma_g$ iff every **recoil** of f is a **descent** of g .

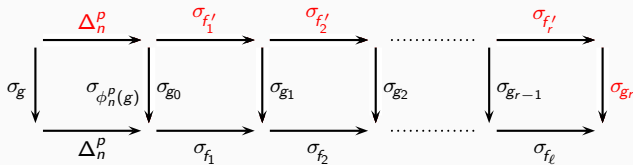
$$i \text{ s.t. } f(i) > f(i+1) \quad i \text{ s.t. } g^{-1}(i) > g^{-1}(i+1)$$

- Proposition (Adjan, El-Rifai–Morton, Thurston, ... 1980s): *Every braid in B_n admits a unique expression $\Delta_n^p \sigma_{f_1} \dots \sigma_{f_r}$ with $p \in \mathbb{Z}$, $f_1 \neq (n, \dots, 2, 1)$, $f_r \neq id$, and every recoil of f_{k+1} is a descent of f_k .*

- The point here: not only theoretical, but also tractable.
 - ▶ The greedy normal form can be computed efficiently.
 - ▶ Key point: computing the normal form of $\sigma_i \beta$ and $\sigma_i^{-1} \beta$ from that of β .

- Recipe:

- ▶ Assume that the normal form of β is $\Delta_n^p \sigma_{f_1} \cdots \sigma_{f_r}$; let σ_g be a permutation-braid;

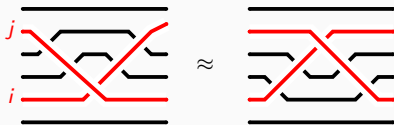


- ▶ The normal form of $\sigma_g \beta$ is $\Delta_n^p \sigma_{f'_1} \cdots \sigma_{f'_p} \sigma_{g_p}$ if $\sigma_{f'_1} \neq \Delta_n$,
and $\Delta_n^{p+1} \sigma_{f'_2} \cdots \sigma_{f'_p} \sigma_{g_p}$ otherwise.

- ▶ And the normal form of $\sigma_g^{-1} \beta$? There exists g' satisfying $\sigma_g \sigma_{g'} = \Delta_n$,
hence $\sigma_g^{-1} = \sigma_{g'} \Delta_n^{-1}$, and $\sigma_g^{-1} \beta = \sigma_{g'} \Delta_n^{p-1} \sigma_{f_1} \cdots \sigma_{f_r}$: continue as above.

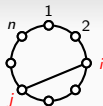
- This corresponds to an **automatic structure** for B_n (Thurston, Cannon),
 - ▶ and, more specifically, to a **Garside structure** (D.–Paris 1997):
 - ↑
 - a submonoid B_n^+ of B_n , plus an element Δ_n of B_n^+ such that
 - B_n is a group of fractions for B_n^+ ,
 - B_n^+ equipped with the (left) divisibility relation is a lattice,
 - $\text{Div}_{\text{left}}(\Delta_n) = \text{Div}_{\text{right}}(\Delta_n)$, $\text{Div}(\Delta_n)$ generates B_n^+ , and $\#\text{Div}(\Delta_n) < \infty$.
 - ▶ Is the Garside structure on B_n unique? Is there another Garside structure on B_n ?
- The **dual** Garside structure on B_n , based on the **Birman–Ko–Lee** generators:

$$\text{for } 1 \leq i < j \leq n: \quad a_{i,j} := \sigma_{j-1} \cdots \sigma_{i+1} \sigma_i \sigma_{i+1}^{-1} \cdots \sigma_{j-1}^{-1}.$$

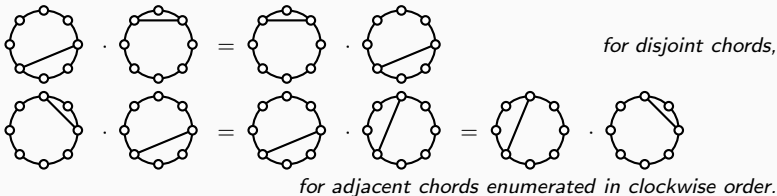


- Definition (Birman–Ko–Lee 1997): B_n^{+*} := submonoid of B_n generated by the $a_{i,j}$'s.
 $\Delta_n^* := a_{1,2} a_{2,3} \cdots a_{n-1,n} (= \sigma_1 \sigma_2 \cdots \sigma_{n-1})$.
- Proposition: (B_n^{+*}, Δ_n^*) is a **Garside structure** on B_n .
 - ▶ a new solution of the Word Problem.

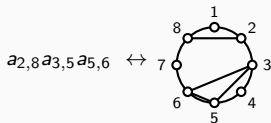
- Chord representation of the Birman–Ko–Lee generators: $a_{i,j} \mapsto$



- Lemma: In terms of the BKL generators, B_n is presented by the relations



- Hence: For P a p -gon, can define a_P to be the product of the $a_{i,j}$ corresponding to $p-1$ adjacent edges of P in clockwise order; *idem* for an union of **disjoint** polygons.



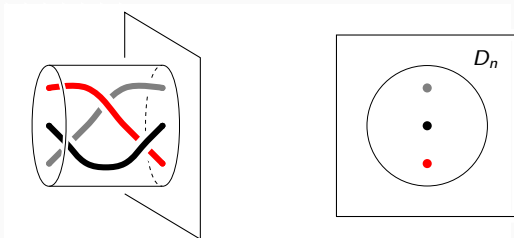
- Proposition (Digne–Michel 2002): The divisors of Δ_n^* in B_n^{+*} are the $\frac{1}{n+1} \binom{2n}{n}$ elements a_P for P a **non-intersecting union of polygons** in an n -punctured circle.

↑
equivalently: a **non-crossing partition** of $\{1, \dots, n\}$

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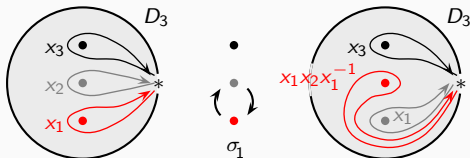
- An n -strand braid diagram = a **danse** of n points in a disk:



... ► an isotopy class of homeomorphisms of D_n leaving ∂D_n fixed
 disk with n marked points \uparrow boundary of D_n

- Proposition: The group B_n is (isomorphic to) the **mapping class group** of D_n .

- Viewing B_n as a group of (isotopy classes of) homeomorphisms of D_n :
 - action of B_n on the fundamental group of D_n , a free group of rank n .



- From there: a homomorphism ρ from B_n to $\text{Aut}(F_n)$:

$$\rho(\sigma_i) : \begin{cases} x_i & \mapsto x_i x_{i+1} x_i^{-1}, \\ x_{i+1} & \mapsto x_i, \\ x_k & \mapsto x_k \text{ for } k \neq i, i+1. \end{cases}$$

- Theorem (Artin): *The homomorphism ρ is injective.*

- a new solution of the Word Problem for B_n (hence of the Braid Isotopy Problem):
a braid word w represents 1 in B_n iff $\rho(w)(x_k) = x_k$ holds for $k = 1, \dots, n$.

- For $x \in \mathbb{Z}$, put $x^+ = \max(0, x)$, $x^- = \min(x, 0)$, and

$$F^+(x_1, y_1, x_2, y_2) = (x_1 + y_1^+ + (y_2^+ - z_1)^+, y_2 - z_1^+, x_2 + y_2^- + (y_1^- + z_1)^-, y_1 + z_1^+),$$

$$F^-(x_1, y_1, x_2, y_2) = (x_1 - y_1^+ - (y_2^+ + z_2)^+, y_2 + z_2^-, x_2 - y_2^- - (y_1^- - z_2)^-, y_1 - z_2^-),$$
 with $z_1 = x_1 - y_1^- - x_2 + y_2^+$ and $z_2 = x_1 + y_1^- - x_2 - y_2^+$.

- Define an **action** of n -strand braid words on \mathbb{Z}^{2n} by

$$(a_1, b_1, \dots, a_n, b_n) * \sigma_i^e = (a'_1, b'_1, \dots, a'_n, b'_n)$$

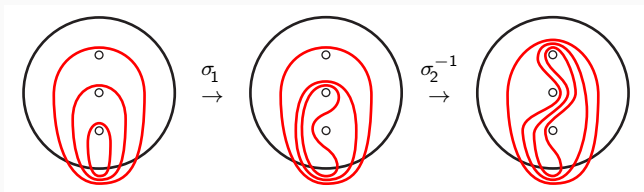
with $a'_k = a_k$ and $b'_k = b_k$ for $k \neq i, i+1$, and $(a'_i, b'_i, a'_{i+1}, b'_{i+1}) = F^e(a_i, b_i, a_{i+1}, b_{i+1})$.

- Definition: The **coordinates** of an n -strand braid word w are $(0, 1, 0, 1, \dots, 0, 1) * w$.

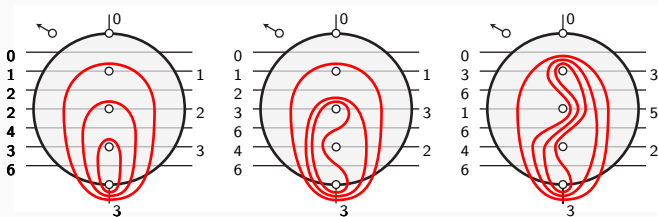
Theorem (Dybnikov 2000): *The coordinates of w only depend on the braid represented by w , and they **characterize** the latter.*

- ▶ Hence: a new solution of the Braid Isotopy Problem:
a braid word w represents 1 iff its Dybnikov coordinates are $(0, 1, 0, 1, \dots, 0, 1)$.
- ▶ An **extremely** efficient method: “**linear** space, **quadratic** time complexity”

- Braid = homeomorphism of D_n ▶ acts on curves drawn in D_n .



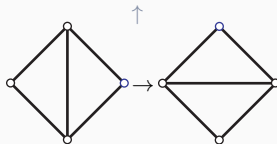
- Count intersections with a fixed triangulation:



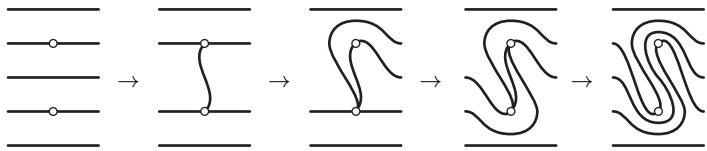
- ▶ $3n + 3$ numbers, which determine the braid

- Fact: The Dynnikov coordinates are the half-differences between the previous intersection numbers.
(going from $3n + 3$ down to $2n$)
- Problem: Compute the coordinates of $\beta\sigma_i^{\pm 1}$ from those of β and i .
 - ▶ compare the intersections of L and $\sigma_i(L)$ with the (fixed) triangulation T
- Main observation:

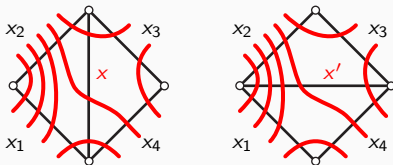
$$\#(\sigma_i(L) \cap T) = \#(L \cap \sigma_i^{-1}(T)).$$
 - ▶ compare the intersections of L with T and $\sigma_i^{-1}(T)$.
- Lemma: If T, T' are any two (singular) triangulations, one can go from T to T' using a finite sequence of *flips*.



- Hence: One **must** go from T to $\sigma_i^{-1}(T)$ by a finite sequence of flips.



- For **one** flip, the formula is



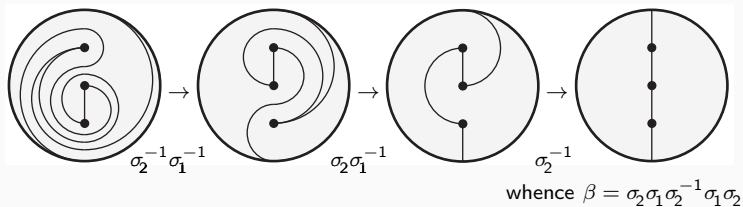
$$x + x' = \max(x_1 + x_3, x_2 + x_4)$$

- Dynnikov's formulas when iterating four times (four flips).

Plan:

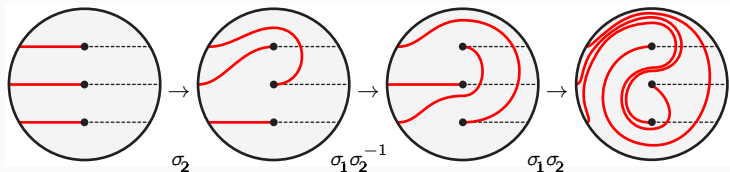
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- Here again: n -strand braid = (isotopy class of) homeomorphism of D_n
- Principle: Fix one (or several) base curve C ,
 - ▶ define a **relaxation strategy** for unbraiding $\beta(C)$ and coming back to C :
 - ▶ the sequence of $\sigma_i^{\pm 1}$ used to unbraid β gives a distinguished expression of β^{-1} (hence a normal form)
 - ▶ requires to define a **complexity** notion first.
- Exemple (Fenn et al. 1997, Dynnikov–Wiest 2006):
 $C =$ **main diameter** of D_n , strategy = consider the “**useful arc**”.



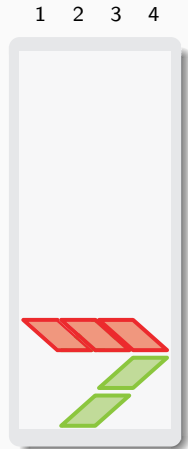
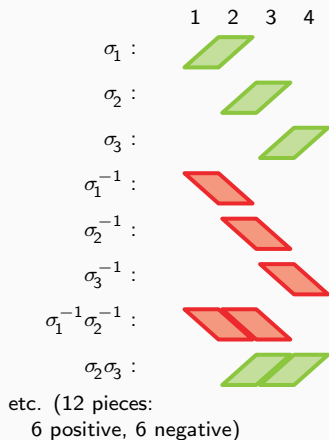
- Exemple 2 (Bressaud 2005):

- ▶ here $C =$ **axes** of standard loops
- ▶ strategy: relax $\beta(x_1)$, then $\beta(x_2), \dots$ by diminishing the number of intersections with half-axes.



- ▶ a normal form on B_n (whence a solution to the Braid Isotopy Problem),
- ▶ together with an **algorithm** computing $NF(w\sigma_i^{\pm 1})$ from $NF(w)$ and i .

- Remark: The Bressaud normal form has nothing to do with positive braids and B_n^+ (nor with B_n^{+*} either).



- **Normal form** of $\sigma_1^{-1}\sigma_2^{-1}\sigma_1\sigma_3^{-1}\sigma_1^{-1}\sigma_2\sigma_3^{-1}\sigma_3$ = $\sigma_2\cdot\sigma_3\cdot\sigma_1^{-1}\sigma_2^{-1}\sigma_3^{-1}$.

On the Garside approach:

- D. Epstein, with J. Cannon, D. Holt, S. Levy, M. Paterson & W. Thurston, *Word Processing in Groups*
Jones & Bartlett Publ. (1992).
- P. Dehornoy, with F. Digne, D. Krammer, J. Michel, *Foundations of Garside Theory*,
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On the Dynnikov coordinates:

- P. Dehornoy, with I. Dynnikov, D. Rolfsen, B. Wiest, *Ordering braids*,
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On relaxation methods:

- R. Fenn, M.T. Greene, D. Rolfsen, C. Rourke, B. Wiest, *Ordering the braid groups*,
Pacific J. of Math. 191 (1999) 49-74.
- X. Bressaud, *A normal form for braids*, J. Knot Th. Ramifications 17-6 (2008) 697-732.

www.math.unicaen.fr/~dehornoy